Verification of GPU Programs: Evaluation Challenges

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PLACES





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Motivation

Static analysis of GPU programs



What is this talk about?

- The previous talk (Liew et al.) presented static analysis of GPU data-races
- In this talk we discuss the challenge of large comparative studies in this area
- We focus on NVidia's CUDA: C++ programs that run on GPU hardware
- In this talk we will refer to GPU programs as kernels

Are there any special requirements for static analysis of kernels?

• Static analysis of kernels requires source code annotations

What verifiers are we comparing?

- Static analysis: Faial (our tool), GPUVerify, PUG
- Also, symbolic execution: GKLEE, SESA



Original program source:

```
--global_-
void saxpy(int n, float a, float *x, float *y)
{
    int i = blockIdx.x*blockDim.x + threadIdx.x;
    if (i < n) y[i] = a*x[i] + y[i];
}</pre>
```

Source annotations for PUG:

```
#include "my_cutil.h"
assume(blockDim.x == 16);
assume(blockDim.y == 16);
assume(gridDim.x == 64);
assume(gridDim.y == 64);
--global___
void saxpy_kernel(int n, float a, float *x, float *y)
{
    int i = blockIdx.x*blockDim.x + threadIdx.x;
    if (i < n) y[i] = a*x[i] + y[i]; }</pre>
```

- PUG requires that a special C header (my_cutil.h) is included
- Static analysis requires problem size (e.g., number of threads) given as assumptions
 - This is the case as SMT solvers cannot handle multiplication of symbolic variables
- PUG expects file extension to be .c rather than the normal .cu
- PUG requires kernel function to end with string kernel



\$ gpuverify --no-inline --only-intra-group --blockDim=16 --gridDim=64 --no-benign-tolerance saxpy-gpuverify.cu

- GPUVerify can accept the number of threads via command-line arguments
- When employing such command-line arguments, GPUVerify can verify some trivial kernels without additional source annotations
 - $\circ~$ This is only the case the only annotation needed is the number of threads
 - Next, we will discuss complex kernels requiring additional annotation
- We implemented our verification tool (Faial) to use GPUVerify style annotations



Source annotations for GPUVerify:

__requires(width = 1024); __requires(height = 1024); __requires(width > gridDim.y); __requires(height > gridDim.x);

```
int xIndex = blockIdx.x * TILE_DIM + threadIdx.x;
int yIndex = blockIdx.y * TILE_DIM + threadIdx.y;
int index = xIndex + width*yIndex;
```

```
for (int r=0; r < nreps; r++) {
    for (int i=0; i<TILE_DIM; i+=BLOCK_ROWS) {
        odata[index+i*width] = idata[index+i*width];
    }
}</pre>
```

Source annotations for PUG:

#include "my_cutil.h"

```
__global__ void kernel (float* odata, float* idata, int width,
        int height, int nreps) {
    assume(width == 1024);
    assume(height = 1024);
    assume(width > gridDim.y);
    assume(height > gridDim.x);
    assume(blockDim.x == 16);
    assume(blockDim.y = 16);
    assume(gridDim.x == 64);
    assume(gridDim.y = 64);
    int xIndex = blockIdx.x * TILE_DIM + threadIdx.x:
    int yIndex = blockIdx.y * TILE_DIM + threadIdx.y;
    int index = xIndex + width*yIndex;
    for (int r=0; r < nreps; r++) {</pre>
        for (int i=0; i<TILE_DIM; i+=BLOCK_ROWS) {</pre>
            odata[index+i*width] = idata[index+i*width]; } }
```

- Additional code annotations are often needed to constrain analysis
- These must be source annotations, and annotation syntax differs by verification tool



Annotation	GPUVerify	PUG
Preconditions	requires();	assume();
File extension	.CU	.C
Number of threads	command-line or source	source annotation
Required headers		"my_cutil.h"

Additionally, PUG places further restrictions on the C++ kernel source code:

- No C++ templates
- No classes
- All loops must be for loops and face additional restrictions

GPUVerify and PUG require differing source annotations!

Source annotations: symbolic execution



Source annotations for GKLEE:

return 0;

```
__global__
void saxpy(int n, float a, float *x, float *y) {
    int i = blockIdx.x*blockDim.x + threadIdx.x;
    if (i < n) y[i] = a*x[i] + y[i]; }</pre>
int main () {
     int n:
     klee_make_symbolic(&n, sizeof(int), "n");
     float a;
     klee_make_symbolic(&a, sizeof(float), "a");
     float *x;
     cudaMalloc((void**)&x, 1024 * sizeof(float));
     float *y;
     cudaMalloc((void**)&y, 1024 * sizeof(float));
     dim3 grid_dim(16);
     dim3 block_dim(64);
     saxpy <<< grid_dim, block_dim >>>(n, a, x, y);
```

Symbolic execution annotations:

- A main function is required to be an execution entry point
- Each kernel parameter must be initialized (symbolically)
- The problem size (e.g., number of threads) must be specified
- Finally, the main function must invoke the kernel function
- Note that this main function specific to symbolic evaluation
 - This main function won't work for running the kernel on a GPU

Related work



- Studies for the following verifiers compare two or fewer total verifiers:
 - 2 for GPUVerify (2012 by Betts et al.)
 - 1 for PUG (2012 by Li and Gopalakrishnan)
 - 2 for GKLEE (2012 by Guodong et al.)
 - 2 for SESA (2014 by Li, Li, and Gopalakrishnan)
- Studies for the following verifiers have limited average lines of code (LoC) analyzed:
 - 13 avg. LoC for ESBMC-GPU (2016 by Pereira et al.)
 - $\circ~16~avg.$ LoC for Simulee (2020 by Wu et al.)

We survived existing published studies: they lack tool diversity and depth!

The challenge of large evaluations

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- Each verifier requires tool-spesific code annotations
- These code annotations are incompatable across verifiers
- Dataset must maintain variations of each program for each verifier

This maintance would be highly labor intensive for researchers!

Source annotations are an impedement to large evaluations

Our evaluation framework

A tool-agnostic kernel format



```
grid_dim = [ 64 ]
block_dim = [ 16 ]
pass = true
body =
       .....
  int i = blockIdx.x*blockDim.x
          + threadIdx.x;
 if (i < n) y[i] = a*x[i] + y[i];
111
scalars = [{n = "int"}, {a = "float"}]
arrays = [{x = "float"}, {y = "float"}]
```

Our tool-agnostic format can specify:

- Number of threads (block_dim, grid_dim)
- C headers to be included
- Kernel function parameters, by type
- Wether the test passes (e.g., data-race freedom, barrier divergence, etc)
- Preamble code
- Preconditions (e.g., invariants)
- The body of the kernel function

Language) to structure this data

• Why TOML? TOML libraries exist to be used with most programming languages

Automation



We have built an ecosystem of tools on top of this tool-agnostic kernel format. Examples include:

- Verifier-specific kernel generation
- Running multiple verifiers on a kernel
- Gathering metrics on code features for each kernel
- Conversion from CUDA to tool-agnostic kernel format

We will cover each of these tasks in the following slides

Verifier-specific kernel generation



Snippet from the GPUVerify template that generates preconditions:

```
{% if pre | length > 0 %}
    /* kernel pre-conditions */
    {% for p in pre %}
    __requires({{ p }});
    {% endfor %}
{% endif %}
```

- We employ the Jinja template engine and Python to generate programs
- Jinja (programming language agnostic) is commonly used to generate HTML web pages

We developed a tool kernel-gen.py to generate kernels:

- Input is the tool-agnostic format
- A command-line argument specifies the format of the output program
- Output is a CUDA kernel formatted for the specified verifier



Running multiple verifiers on a kernel

\$ <mark>test</mark> -tools.py saxpy.toml RUN timeout 60 faialparse-gv-args saxpy-faial.cu						
RUN timeout 60 gpuverifyno-inlineonly-intra-group <mark>blockDim=16gridDim=64</mark>						
no-benign-tolerance saxpy-gpuverify.cu						
RUN timeout 60 pug saxpy-pug.c						
status	time	memory	data_races	tool		
0	0.18	48.418	drf	faial		
0	1.34	37.5859	drf	gpuverify		
0	0.05	50.6484	drf	pug		

- Some verifiers require metadata stored in the tool-agnostic kernel format
- For example, recall that GPUVerify needs the number of threads

Gathering metrics on code features



We extend our tool-agnostic kernel format to store metadata on code features of the dataset. The following is an example of such metadata:

```
max_sync_nesting = 1
sync_loop_count = 1
unsync_loop_count = 1
loop_count = 2
write_count = 3
read_count = 4
if_count = 0
sync_count = 2
line_count = 71
```

This is possible as the TOML format is easily extensible

Conversion from CUDA to agnostic format

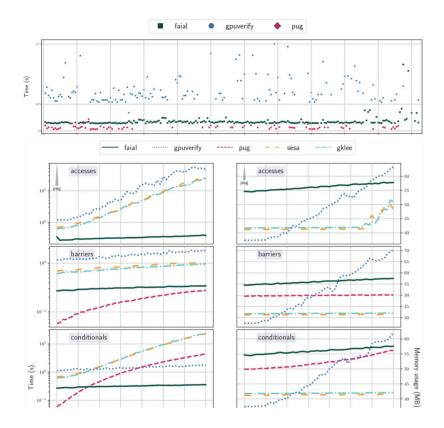


- Much of our evaluation is based on on GPUVerify's CAV 2014 dataset
- This dataset of 227 CUDA kernels comes with GPUVerify source annotations
- We employed a script to convert this CUDA dataset to our tool-agnostic format
- The script (~150 lines of Python) consists of rudimentary text processing
- This conversion was mostly-automated, handling 65% of the dataset

Takeaway: a little automation saves a lot of time!

Results: CAV 2021 evaluation





This benchmarking infrastructure enabled multiple experiments:

- A dataset of 227 CUDA kernels, 3 static verifier, including kernels with up to 850 lines of code
- And an additional 250 synthetic kernels, 5 static and symbolic verifiers
- To our knowledge, this is the largest published experimental evaluation of GPU verifiers to date



Conclusion

Future challenges



Support for numeric and symbolic problems

- Currently, we check verifier output for boolean correctness tests
- In the future, we are interested in parsing output for numeric and symbolic output

Weakening and strengthening preconditions

• We are interested in which preconditions affect certain properties

What we learned



- Source annotations are an impediment to large evaluations
- Templates aid the handling of differing annotations via program generation
- A little automation goes a long way

With templates and automation in place, larger evaluations are feasible.